



10. Network planning and dimensioning

lect10.ppt

S-38.145 - Introduction to Teletraffic Theory - Fall 1999

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10. Network planning and dimensioning

Literature

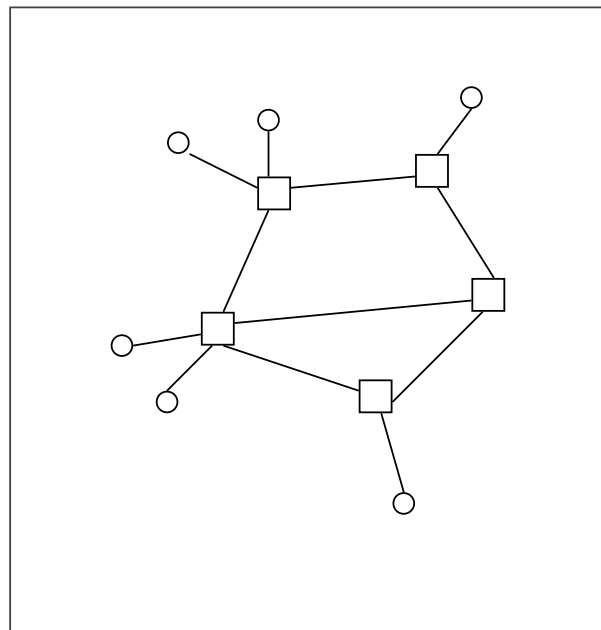
- 1 A. Olsson, ed. (1997)
 - “Understanding Telecommunications 1”
 - Studentlitteratur, Lund, Sweden
- 2 A. Girard (1990)
 - “Routing and Dimensioning in Circuit-Switched Networks”
 - Addison-Wesley, Reading, MA

Contents

- Introduction
- Network planning
- Traffic forecasts
- Dimensioning

Telecommunication network

- A simple model of a telecommunication network consists of
 - **nodes**
 - terminals ○
 - network nodes □
 - **links** between nodes
- **Access network**
 - connects the terminals to the network nodes
- **Trunk network**
 - connects the network nodes to each other



Why network planning and dimensioning?

- “The purpose of dimensioning of a telecommunications network is to ensure that

the expected needs will be met in an economical way

both for subscribers and operators.”

Source: [1]

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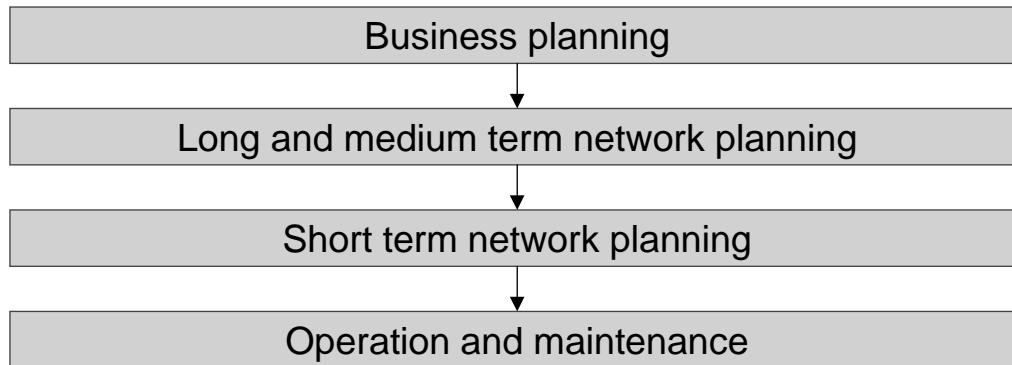
Contents

- Introduction
- Network planning
- Traffic forecasts
- Traffic dimensioning

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Network planning in a stable environment (1)

- Traditional planning situation:



Source: [1]

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Network planning in a stable environment (2)

- Traffic aspects
 - Data collection (current status)
 - traffic measurements
 - subscriber amounts and distribution
 - Forecasting
 - service scenarios
 - traffic volumes and profiles
- Economical aspects
- Technical aspects
- Network optimisation and dimensioning
 - hierarchical structure and topology
 - traffic routing and dimensioning
 - circuit routing

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Traditional planning process by Girard (1)

- As with any decision process, network planning relies on **external information**
 - Forecast of demand for services over some planning horizon
 - Economic information concerning the cost structure of the network elements and maintenance
 - Knowledge about the technical capabilities of the available systems
- The planning problem can now be stated as follows:
 - to implement the first four layers of the OSI model
 - to provide the required physical support

Source: [2]

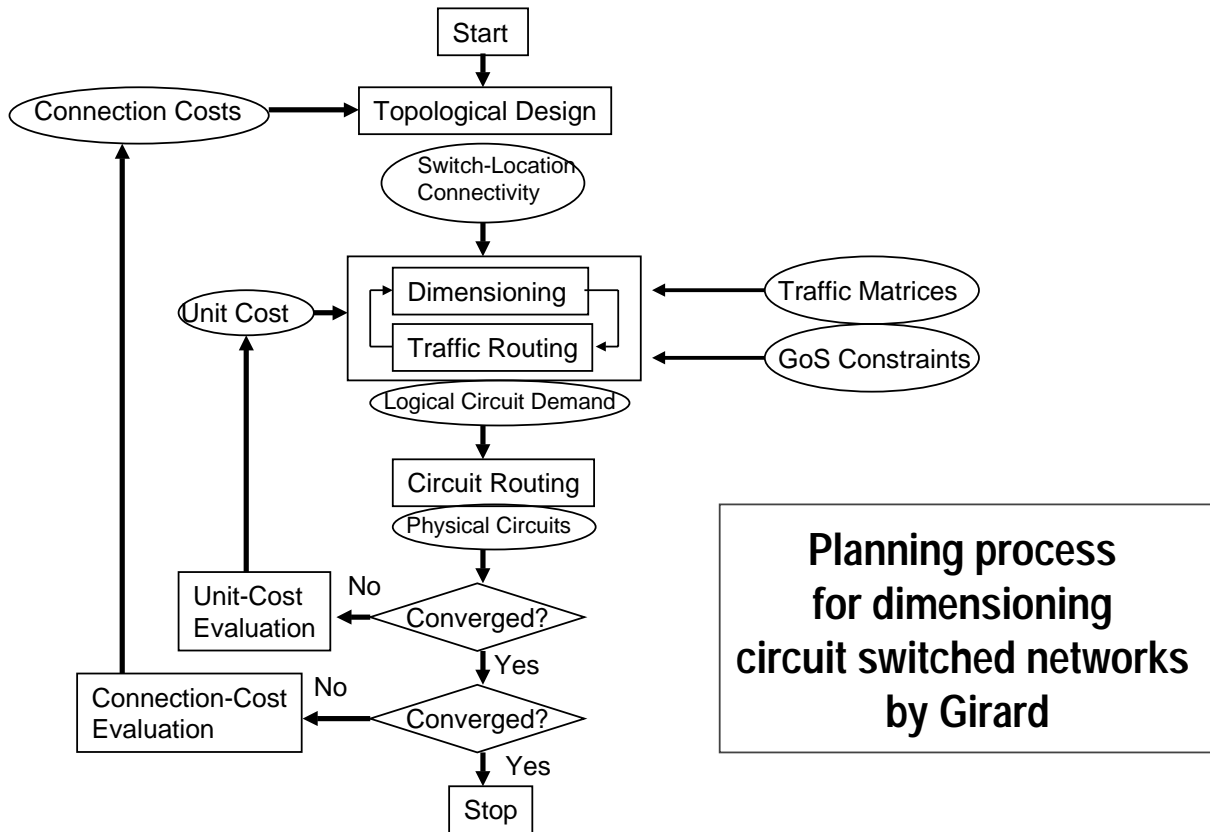
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Traditional planning process by Girard (2)

- Assuming that all the protocol issues have been settled and the transmission technology is known, what remains is a complex, distributed and dynamic capacity-augmentation problem
 - only feasible solution approach: decomposition and iteration
- Stages of the planning process:
 - Topological design
 - Network-synthesis problem
 - Traffic routing
 - Dimensioning
 - Network-realization (circuit-routing) problem
- These four stages are interrelated
 - ⇒ the planning process is **iterative** (at many levels)
- Different **planning horizons** at various stages

Source: [2]

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**Planning process
for dimensioning
circuit switched networks
by Girard**

Source: [2]

Traditional planning process by Girard (3)

- Topological design:
 - Determine where to place components and how to interconnect them
 - By methods of **topological optimization** and **graph theory**
 - Input:
 - information about transmission network summarized into a fixed interconnection cost per unit length between offices
 - switch costs depending just on the switching technology
 - Output:
 - connectivity matrix
 - optimal location of switches or concentrators (optionally)

Source: [2]

Traditional planning process by Girard (4)

- Network synthesis:
 - Calculate the optimal size of the components (that is: the transmission and switching systems) within the topology specified and subject to GoS constraints on network-performance measures
 - By methods of **nonlinear optimization**
 - Input:
 - topology, traffic matrices, GoS constraints, cost function (unit cost)
 - Output:
 - route plan
 - set of logical links between the nodes
(that is: requirements for transmission facilities betw. switching points)
 - Comprises of two iterated substages:
 - Traffic routing
 - Dimensioning
 - Specific to telecommunications!

Source: [2]

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Traditional planning process by Girard (5)

- Traffic routing:
 - Determine how to connect calls as they arrive, given the topology and size of the components
- Dimensioning:
 - Determine the size of the components subject to GoS constraints and given the topology and a routing method

Source: [2]

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Traditional planning process by Girard (6)

- Network realization:
 - Determine how to implement the capacity requirement (for transmission and switching equipments) using the available components and taking further into account **reliability** (\Rightarrow multipath routing)
 - By methods of **multicommodity flow optimization**
 - Input:
 - logical-circuit demand
 - fixed costs, module costs and reliability of available components
 - other reliability requirements
 - Output:
 - physical circuits plan
 - detailed information of actual transmission cost between nodes

Source: [2]

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Network planning in a turbulent environment (1)

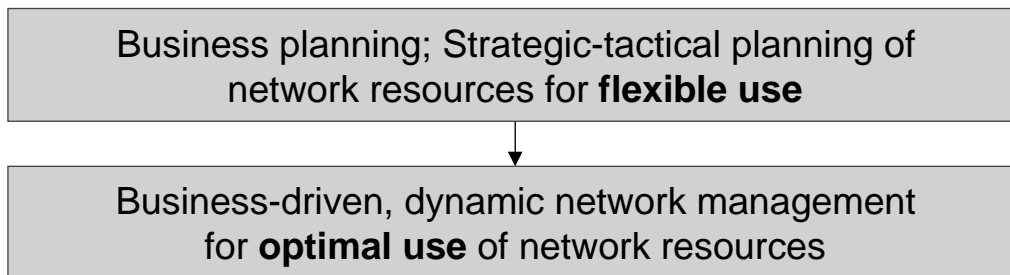
- Additional decision data are needed from the following areas:
 - The market, with regard to a specific business concept
 - due to competition!
 - operator's future role (niche): dominance/co-operation
 - Customer demands:
 - new services: Internet & mobility (first of all)
 - new business opportunities
 - Technology:
 - new technology: ATM, xDSL, GSM, CDMA, WDM
 - Standards:
 - new standards issued continuously
 - Operations and network planning support:
 - new computer-aided means
 - Costs:
 - trends: equipment costs going down, staff costs going up

Source: [1]

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Network planning in a turbulent environment (2)

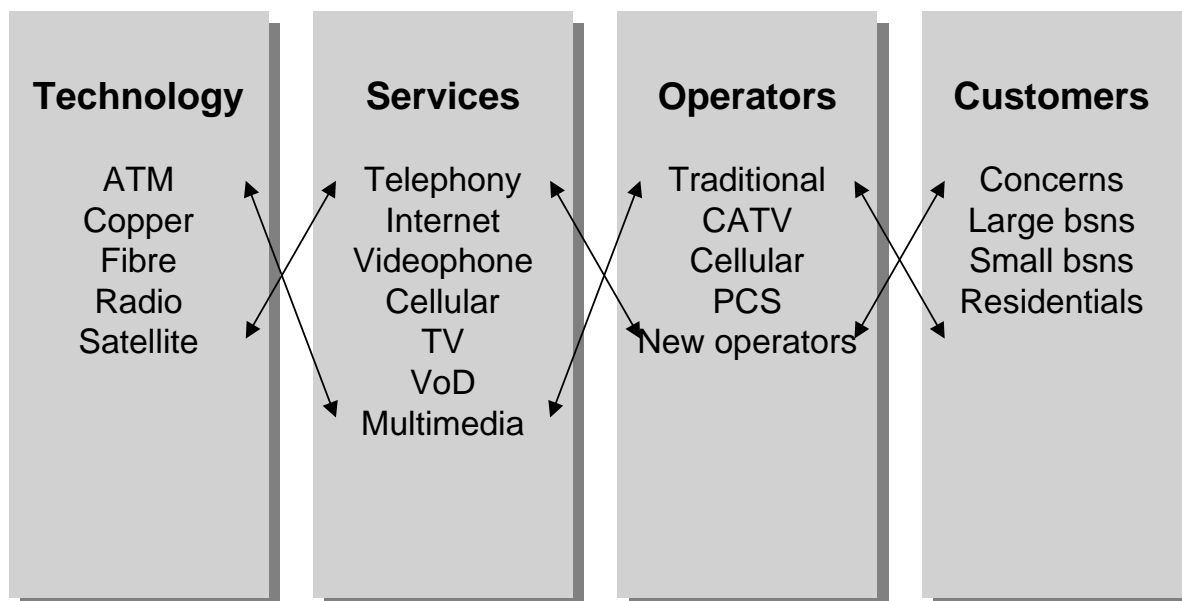
- Safeguards for the operator:
 - Change the network architecture so that it will be more **open**, with generic **platforms**, if possible
 - Build the network with a certain prognosticated overcapacity (**redundancy**) in generic parts where the marginal costs are low
- New planning situation (shift of focus to a strategic-tactical approach):



Source: [1]

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"The new conception of the world"



Source: [1]

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Contents

- Introduction
- Network planning
- **Traffic forecasts**
- Traffic dimensioning

Need for traffic measurements and forecasts

- To properly dimension the network we need to
estimate the traffic offered
- If the network is already operating,
 - the current traffic is most precisely estimated by making traffic measurements
- Otherwise, the estimation should be based on other information, e.g.
 - estimations on characteristic traffic generated by a subscriber
 - estimations on the number of subscribers
- Long time-span of network investments \Rightarrow
 - it is not enough to estimate only the current traffic
 - forecasts of future traffic are also needed

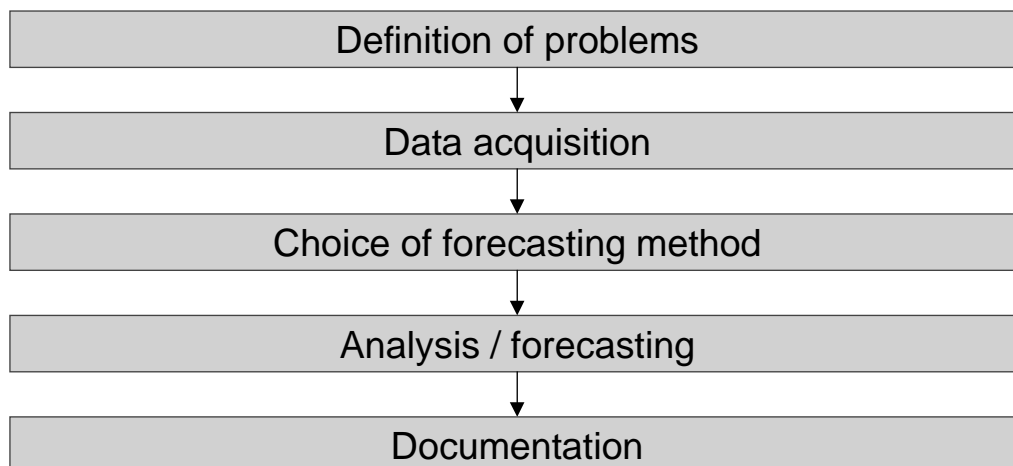
Traffic forecasting

- Information about future demands for telecommunications
 - an estimation of future tendency or direction
- Purpose
 - provide a basis for decisions on investments in network
- Forecast periods
 - time aspect important (reliability)
 - need for forecast periods of different lengths

Source: [1]

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Forecasting procedure



Source: [1]

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Forecasting methods

- Trend methods
 - linear extrapolation
 - nr of subscribers increased yearly by about 200 in the past 5 years
⇒ $3 * 200 = 600$ new subscribers in the next 3-year period
 - not suitable if growth is exponential
- Statistical demand analysis
 - network operator seeks to map out those factors that underlie the earlier development
 - changes that can be expected during the forecasting period are then collated
- Assessment methods
 - analogy method: situations or objects with similar preconditions will develop similarly

Source: [1]

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Traffic forecast

- Traffic forecast defines
 - the estimated traffic growth in the network over the planning period
- Starting point:
 - current traffic volume during busy hour (measured/estimated)
- Other affecting factors:
 - changes in the number of subscribers
 - change in traffic per subscriber (characteristic traffic)
- Final result (that is, the forecast):
 - **traffic matrix** describing the traffic **interest** between exchanges (traffic areas)

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Traffic matrix

- The final result of the traffic forecast is given by a **traffic matrix**
- Traffic matrix $T = (T(i,j))$
 - describes traffic interest between exchanges
 - N^2 elements ($N = \text{nr of exchanges}$)
 - element $T(i,i)$ tells the estimated traffic within exchange i
 - element $T(i,j)$ tells the estimated traffic from exchange i to exchange j
- Problem:
 - easily grows too big: 600 exchanges \Rightarrow 360,000 elements!
- Solution: hierarchical representation
 - higher level: traffic between traffic areas
 - lower level: traffic between exchanges within one traffic area

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Example (1)

- **Data:**
 - There are 1000 private subscribers and 10 companies with their own PBX's in the area of a local exchange.
 - The characteristic traffic generated by a private subscriber and a company are estimated to be 0.025 erlang and 0.200 erlang, respectively.
- **Questions:**
 - What is the total traffic intensity a generated by all these subscribers?
 - What is the call arrival rate λ assumed that the mean holding time is 3 minutes?
- **Answers:**
 - $a = 1000 * 0.025 + 10 * 0.200 = 25 + 2 = \mathbf{27 \text{ erlangs}}$
 - $h = 3 \text{ min}$
 - $\lambda = a/h = 27/3 \text{ calls/min} = \mathbf{9 \text{ calls/min}}$

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Example (2)

- **Data:**
 - In a 5-year forecasting period the number of new subscribers is estimated to grow linearly with rate 100 subscribers/year.
 - The characteristic traffic generated by a private subscriber is assumed to grow to value 0.040 erlang.
 - The total nr of companies with their own PBX is estimated to be 20 at the end of the forecasting period.
- **Question:**
 - What is the estimated total traffic intensity a at the end of the forecasting period?
- **Answer:**
 - $a = (1000 + 5 \cdot 100) \cdot 0.040 + 20 \cdot 0.200 = 60 + 4 = \mathbf{64 \text{ erlangs}}$

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Example (3)

- **Data:**
 - Assume that there are three similar local exchanges.
 - Assume further that one half of the traffic generated by a local exchange is local traffic and the other half is directed uniformly to the two other exchanges.
- **Question:**
 - Construct the traffic matrix T describing the traffic interest between the exchanges at the end of the forecasting period.

- **Answer:**
 - $T(i,i) = 64/2 = 32 \text{ erlangs}$
 - $T(i,j) = 64/4 = 16 \text{ erlangs}$

area	1	2	3	sum
1	32	16	16	64
2	16	32	16	64
3	16	16	32	64
sum	64	64	64	192

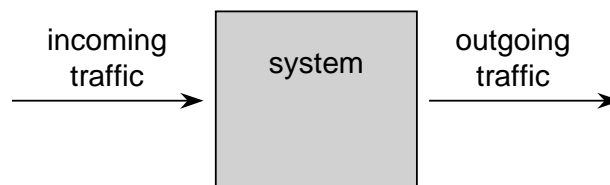
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Traffic dimensioning (1)

- Telecommunications system from the traffic point of view:



- Basic task in **traffic dimensioning**:

Determine the minimum **system capacity** needed in order that the incoming **traffic** meet the specified **grade of service**

Traffic dimensioning (2)

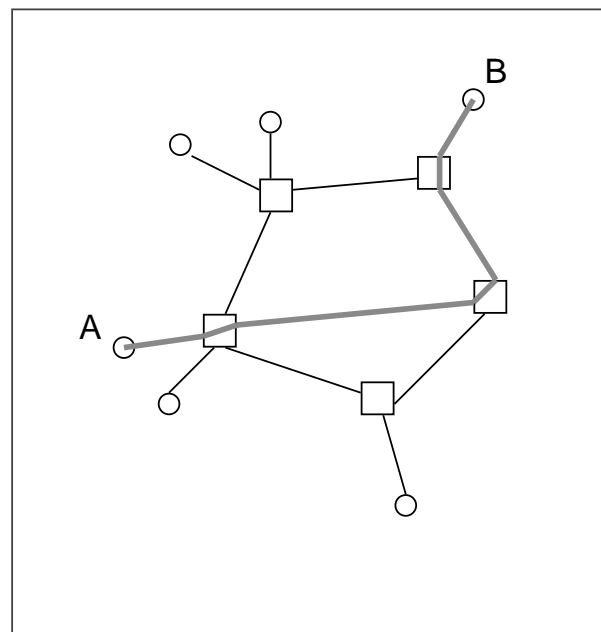
- Observation:
 - Traffic is varying in time
- General rule:
 - Dimensioning should be based on peak traffic not on average traffic
- However,
 - Revenues are based on average traffic
- For dimensioning (of telephone networks), peak traffic is defined via the concept of busy hour:

Busy hour \approx the continuous 1-hour period for which the traffic volume is greatest

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Telephone network model

- Simple model of a telephone network consists of
 - network nodes (exchanges)
 - links between nodes
- Traffic consists of **calls**
- Each call has two phases
 - first, the connection has to set up through the network (**call establishment** phase)
 - only after that, the information transfer is possible (**information transfer** phase)



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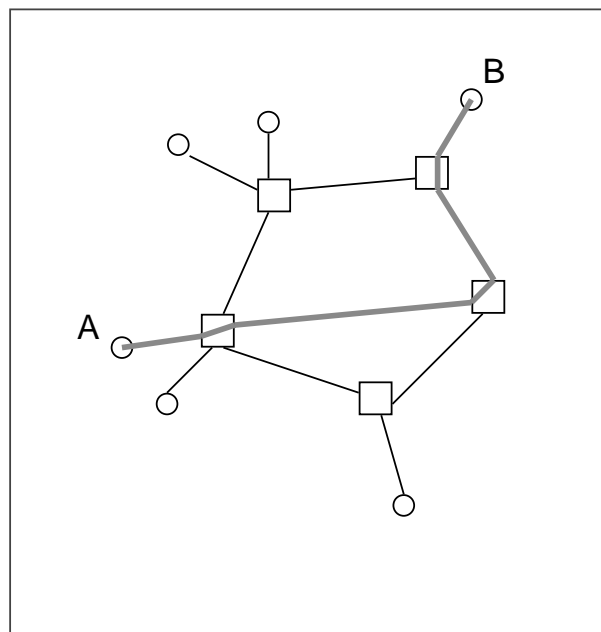
Two kinds of traffic processes

- Traffic process in each **network node**
 - due to **call establishments**
 - during the call establishment phase
 - each call needs (and competes for) processing resources in each network node (switch) along its route
 - it typically takes **some seconds** (during which the call is processed in the switches, say, **some milliseconds**)
- Traffic process in each **link**
 - due to **information transfer**
 - during the information transfer phase
 - each call occupies one channel on each link along its route
 - information transfer lasts as long as one of the participants disconnects
 - ordinary telephone calls typically hold **some minutes**
- **Note:** totally **different time scales** of the two processes

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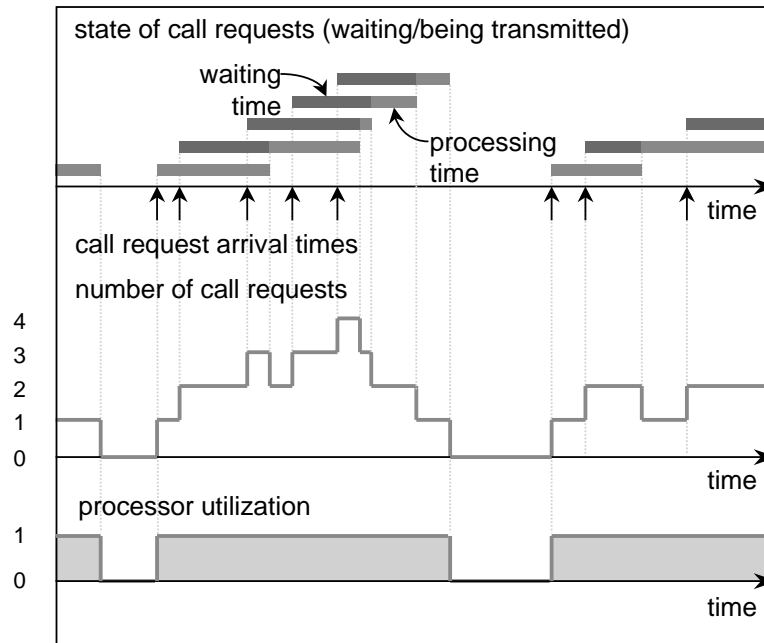
Simplified traffic dimensioning in a telephone network

- Assume
 - fixed topology and routing
 - given traffic matrix
 - given GoS requirements
- Dimensioning of network nodes: Determine the required **call handling capacity**
 - max number of call establishments the node can handle in a time unit
- Dimensioning of links: Determine the required **number of channels**
 - max number of ongoing calls on the link



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Traffic process during call establishment (1)



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Traffic process during call establishment (2)

- Call (request) arrival process is modelled as
 - a Poisson process with intensity λ
- Further we assume that call processing times are
 - IID and exponentially distributed with mean s
 - typically s is in the range of **milliseconds** (not minutes as h)
 - s is more a **system parameter** than a traffic parameter
- Finally we assume that the call requests are processed by
 - a single processor with an infinite buffer
- The resulting traffic process model is
 - the **M/M/1 queueing model** with traffic load $\rho = \lambda s$

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Traffic process during call establishment (3)

- Pure delay system \Rightarrow

Grade of Service measure = Mean waiting time $E[W]$

- Formula for the mean waiting time $E[W]$ (assuming that $\rho < 1$):

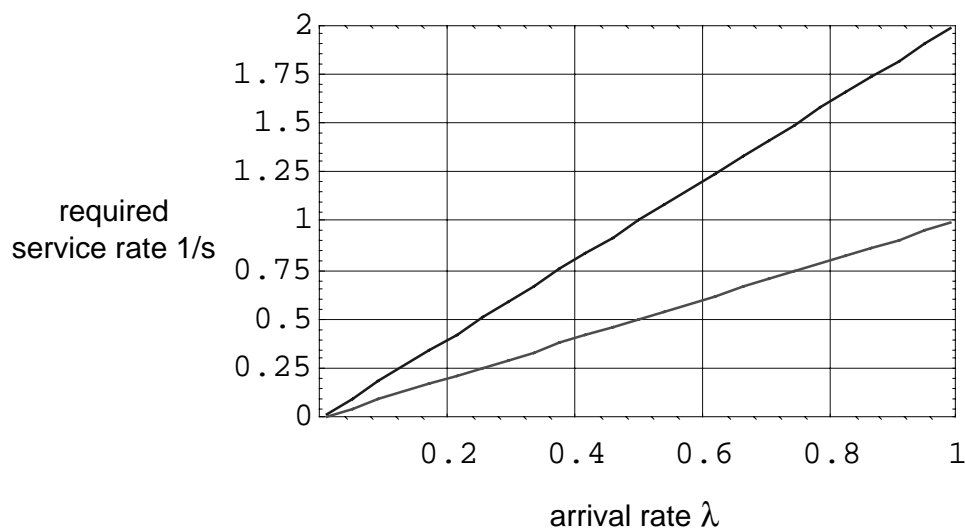
$$E[W] = s \cdot \frac{\rho}{1-\rho}$$

- $\rho = \lambda s$
- **Note:** $E[W]$ grows to infinity as ρ tends to 1

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Dimensioning curve

- Grade of Service requirement: $E[W] \leq s$
 \Rightarrow Allowed load $\rho \leq 0.5 = 50\% \Rightarrow \lambda s \leq 0.5$
 \Rightarrow Required service rate $1/s \geq 2\lambda$



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Dimensioning rule

- To get the required Grade of Service (the average time a customer waits before service should be less than the average service time) ...

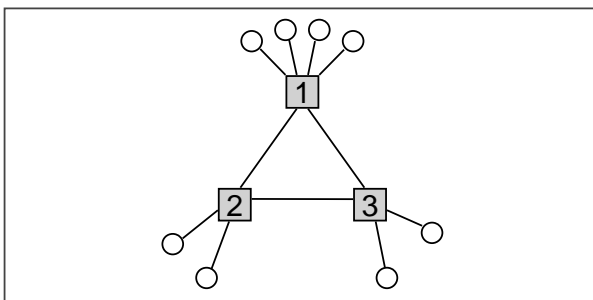
... Keep the traffic load less than 50%

- If you want a less stringent requirement, still remember the **safety margin** ...

Don't let the total traffic load approach to 100%

- Otherwise you'll see an explosion!

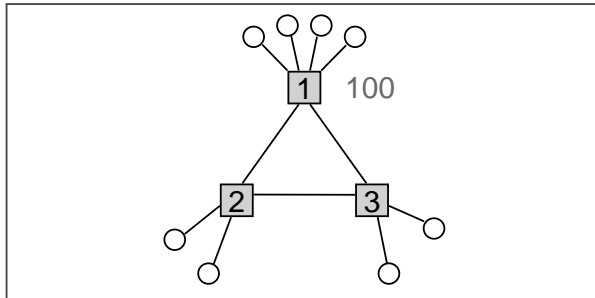
Example (1)



area	1	2	3	sum
1	60	15	15	90
2	30	30	15	75
3	30	15	30	75
sum	120	60	60	240

- Assumptions:**
 - 3 local exchanges completely connected to each other
 - Traffic matrix T describing the busy hour traffic interest (in erlangs) given below
 - Fixed (direct) routing: calls are routed along shortest paths.
 - Mean holding time $h = 3$ min.
- Task:**
 - Determine the call handling capacity needed in different network nodes according to the GoS requirement $\rho < 50\%$

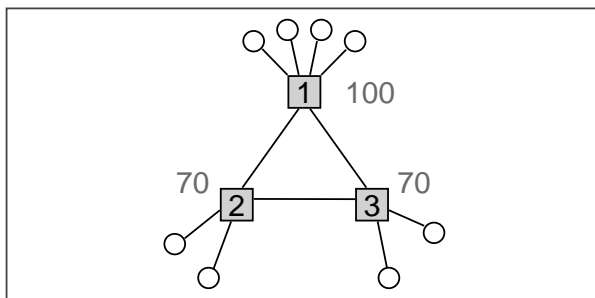
Example (2)



area	1	2	3	sum
1	60	15	15	90
2	30	30	15	75
3	30	15	30	75
sum	120	60	60	240

- **Node 1:**
 - call requests from own area:
 $[T(1,1) + T(1,2) + T(1,3)]/h = 90/3 = 30$ calls/min
 - call requests from area 2:
 $T(2,1)/h = 30/3 = 10$ calls/min
 - call requests from area 3:
 $T(3,1)/h = 30/3 = 10$ calls/min
 - total call request arrival rate:
 $\lambda(1) = 30+10+10 = 50$ calls/min
 - required call handling capacity:
 $\rho(1) = \lambda(1)/\mu(1) = 0.5 \Rightarrow \mu(1) = 2*\lambda(1) = \mathbf{100}$ calls/min

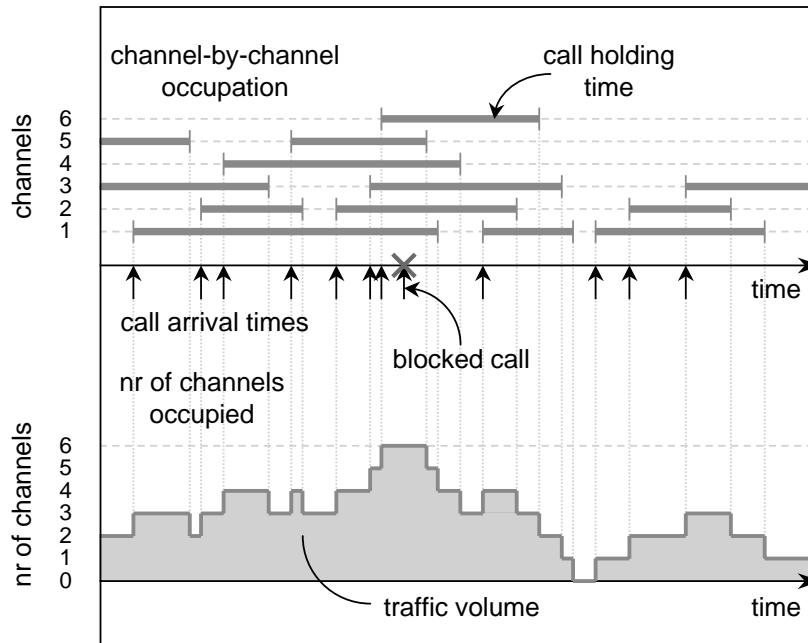
Example (3)



area	1	2	3	sum
1	60	15	15	90
2	30	30	15	75
3	30	15	30	75
sum	120	60	60	240

- **Node 2:**
 - total call request arrival rate:
 $\lambda(2) = [T(2,1) + T(2,2) + T(2,3) + T(1,2)+T(3,2)]/h = (75+15+15)/3 = 35$ calls/min
 - required call handling capacity:
 $\mu(2) = 2*\lambda(2) = \mathbf{70}$ calls/min
- **Node 3:**
 - total call request arrival rate :
 $\lambda(3) = [T(3,1) + T(3,2) + T(3,3) + T(1,3)+T(2,3)]/h = (75+15+15)/3 = 35$ calls/min
 - required call handling capacity:
 $\mu(3) = 2*\lambda(3) = \mathbf{70}$ calls/min

Traffic process during information transfer (1)



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Traffic process during information transfer (2)

- Call arrival process has already been modelled as
 - a Poisson process with intensity λ
- Further we assume that call holding times are
 - IID and generally distributed with mean h
 - typically h is in the range of **minutes** (not milliseconds as s)
 - h is more a **traffic parameter** than a system parameter
- The resulting traffic process model is
 - the **M/G/n/n loss model** with (offered) traffic intensity $a = \lambda h$

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Traffic process during information transfer (3)

- Pure loss system \Rightarrow

Grade of Service measure = Call blocking probability B

- Erlang's blocking formula:

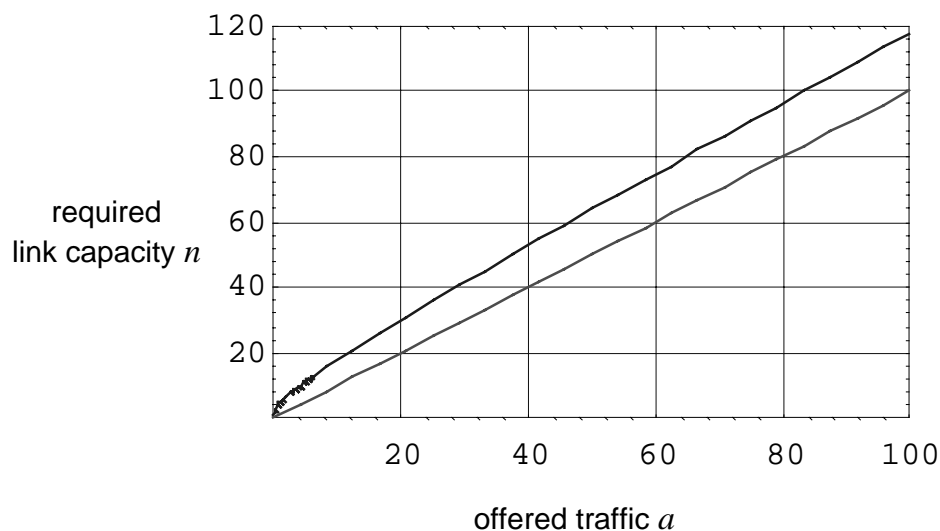
$$B = \text{Erl}(n, a) = \frac{\frac{a^n}{n!}}{\sum_{i=0}^n \frac{a^i}{i!}}$$

- $a = \lambda h$
- $n! = n(n-1)(n-2) \dots 1$

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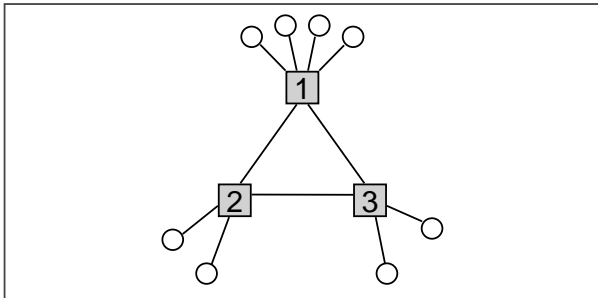
Dimensioning curve

- Grade of Service requirement: $B \leq 1\%$
 \Rightarrow Required link capacity: $n = \min\{i = 1, 2, \dots \mid \text{Erl}(i, a) \leq B\}$



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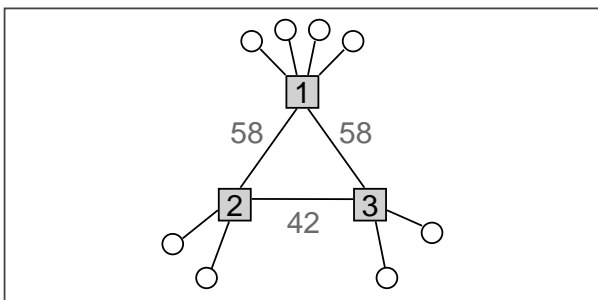
Example (1)



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3	30	15	30	75
sum	120	60	60	240

- **Assumptions:**
 - 3 local exchanges completely connected to each other with two-way links
 - Traffic matrix T describing the busy hour traffic interest (in erlangs) given below
 - Fixed (direct) routing: calls are routed along shortest paths.
 - Mean holding time $h = 3$ min.
- **Task:**
 - Dimension trunk network links according to the GoS requirement $B < 1\%$

Example (2)



area	1	2	3	sum
1	60	15	15	90
2	30	30	15	75
3	30	15	30	75
sum	120	60	60	240

- **Link 1-2 (betw. nodes 1 and 2):**
 - total offered traffic:
 $a(1-2) = T(1,2) + T(2,1)$
 $= 15+30 = 45$ erlang
 - required capacity:
 $n(1-2) = \min\{i \mid \text{Erl}(i,45) < 1\%\}$
 $\Rightarrow n(1-2) = \mathbf{58 \text{ channels}}$
- **Link 1-3:**
 - required capacity:
 $n(1-3) = \min\{i \mid \text{Erl}(i,45) < 1\%\}$
 $\Rightarrow n(1-3) = \mathbf{58 \text{ channels}}$
- **Link 2-3:**
 - required capacity:
 $n(2-3) = \min\{i \mid \text{Erl}(i,30) < 1\%\}$
 $\Rightarrow n(2-3) = \mathbf{42 \text{ channels}}$

Table: $B = \text{Erl}(n,a)$

- $B = 1\%$

n :	a :
– 35 channels	24.64 erlang
– 36 channels	25.51 erlang
– 37 channels	26.38 erlang
– 38 channels	27.26 erlang
– 39 channels	28.13 erlang
– 40 channels	29.01 erlang
– 41 channels	29.89 erlang
– 42 channels	30.78 erlang
– 43 channels	31.66 erlang
– 44 channels	32.55 erlang
– 45 channels	33.44 erlang

- $B = 1\%$

n :	a :
– 50 channels	37.91 erlang
– 51 channels	38.81 erlang
– 52 channels	39.71 erlang
– 53 channels	40.61 erlang
– 54 channels	41.51 erlang
– 55 channels	42.41 erlang
– 56 channels	43.32 erlang
– 57 channels	44.23 erlang
– 58 channels	45.13 erlang
– 59 channels	46.04 erlang
– 60 channels	46.95 erlang

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End-to-end blocking probability

- Thus far we have concentrated on the single link case, when calculating the call blocking probability B_c
- However, there can be many (trunk network) links along the route of a (long distance) call. In this case it is more interesting to calculate the total **end-to-end blocking probability** B_e experienced by the call. A method (called **Product Bound**) to calculate B_e is given below.
- Consider a call traversing through links $j = 1, 2, \dots, J$. Denote by $B_c(j)$ the blocking probability experienced by the call in each single link j . Then

$$B_e = 1 - (1 - B_c(1)) * (1 - B_c(2)) * \dots * (1 - B_c(J))$$

$$B_c(j)\text{'s small} \Rightarrow B_e \approx B_c(1) + B_c(2) + \dots + B_c(J)$$

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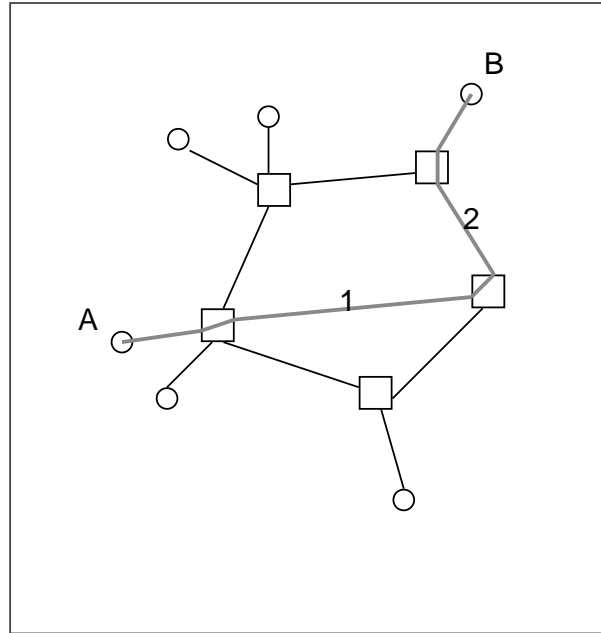
Example

- The call from A to B is traversing through trunk network links 1 and 2
- Let $B_c(1)$ and $B_c(2)$ denote the call blocking probability in these links
- Product Bound (PB):

$$B_e = 1 - (1 - B_c(1)) * (1 - B_c(2))$$

$$= B_c(1) + B_c(2) - B_c(1) * B_c(2)$$
- Approximately:

$$B_e \approx B_c(1) + B_c(2)$$



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THE END



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