



Multi-class routing in DiffServ networks

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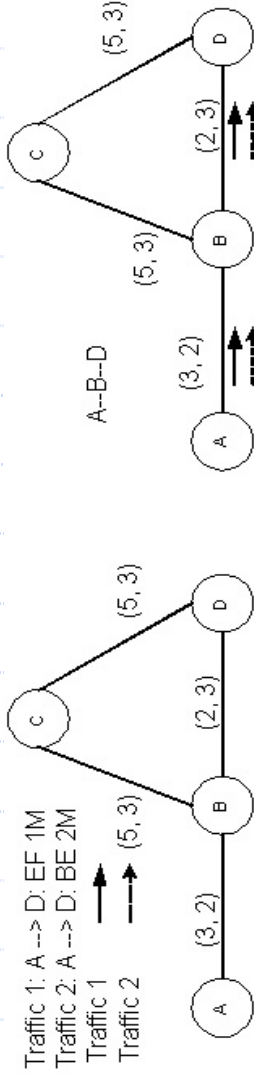
January 8th 2004, IRoNet seminar



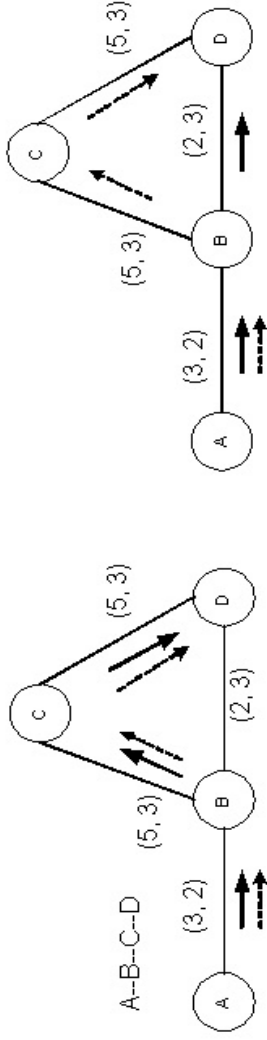
Motivation

- ◆ Best-effort oriented IP networks are expected to migrate into multi-class networks, e.g.,
 - DiffServ networks
 - MPLS networks
- ◆ Multi-class networks are expected to support traffic with
 - Various traffic characteristics, and
 - Various QoS demands
- ◆ Traditional routing belongs to single-class routing (SCR) scenario, which calculates same routing for all kinds of traffic classes.
- ◆ In contrast, we propose multi-class routing (MCR), which calculates different routing for different traffic classes.

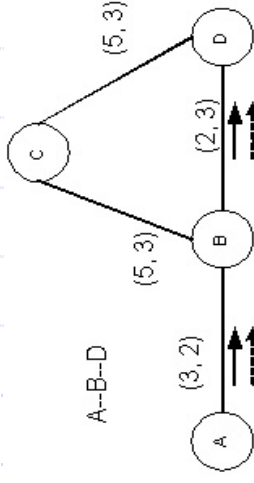
A simple example



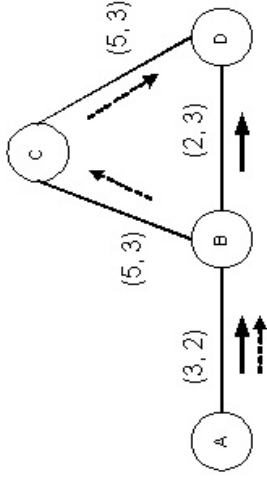
(a) Network and traffic



(c) SCR: BSP



(b) SCR: SP



(d) MCR: (SP, BSP)

◆ Two traffic from A to D

- Traffic 1: EF 1M
- Traffic 2: BE 2M

◆ Goals:

- Minimum delay for EF
- Maximum throughput for BE

◆ A link is denoted by (bandwidth, delay)

Only MCR can achieve routing optimization of each class.

Multi-class routing (MCR)

- ◆ Clarification
 - MCR is related but not same as QoS routing (QoSR)
 - ◆ MCR may use QoS routing algorithms
 - MCR for MPLS is referred to PERD (a legacy name), which is based on requests.
- ◆ We propose MCR mainly because of the following reasons:
 - Different classes may have different QoS requirements, e.g., video/voice traffic conveyed by EF class requires minimum delay and delay jitter while WWW/FTP traffic conveyed by BE class requires maximum throughput.
 - Migrating SCR into MCR requires only a few changes, which is much simpler than MPLS. Memory cost is not significant in modern routers.
 - DiffServ might be a policy based network, where MCR can be used to compute optimal routes that are downloadable to each router.
 - Our previous simulation results prove that MCR can improve the performance of not only high priority classes but also low priority classes.

MCR algorithm

MCR scheme:

```
Begin
Initialization {
  Set routing metric and method for each class,
  i.e., set  $D = \{d_i \mid i \in C\}$ 
  Initialize routes for each class,
  i.e.,  $P = \{P_i \mid i \in C\}$ 
}
Start{
  Step 1:
  Started with the highest priority class, i.e.,  $i=0$ ,
  Compute  $P_i$  with  $d_i$ 
  Step 2:
  Proceed to lower priority class, i.e.,  $i++$ ;
  Compute  $P_i$  with  $d_i$ 
  Step 3:
  If this is not the lowest priority class,
  go to step 2.
}
End.
```

Notations:

$C = \{i \mid i = 0, 1, 2, \dots\}$, the set of traffic classes. These classes are ordered from the highest priority to the lowest priority.

$M = \{m_j \mid j = 0, 1, 2, \dots\}$, the set of routing metrics.

$R = \{r_k \mid k = 0, 1, 2, \dots\}$, the set of routing methods.

$d_i = \{m_j, r_k \mid m_j \in M, r_k \in R\}$, a tuple of routing metric and routing method for traffic class i .

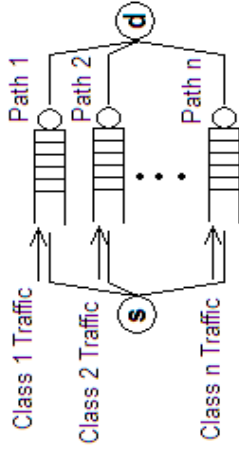
P_i , the set of routes for traffic class i .

MCR modeling

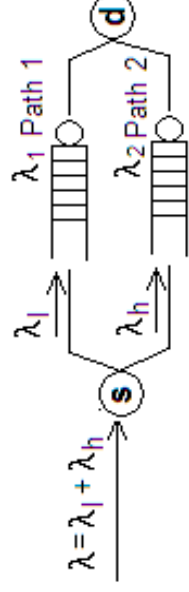
Class 1 Traffic + Class 2 Traffic + ... + Class n Traffic



(a) SCR



(b) MCR



$$\begin{aligned} \Delta T &= T_{P1}^{MCR} - T_{P1}^{SP} \\ &= \frac{\lambda_l + \lambda_h}{\mu_{P1}(\mu_{P1} - \lambda_l - \lambda_h)} - \frac{\lambda_l + \lambda}{\mu_{P1}(\mu_{P1} - \lambda_l - \lambda)} \\ &= - \frac{\lambda_h}{(\mu_{P1} - \lambda_l - \lambda_h)(\mu_{P1} - \lambda_l - \lambda)} \end{aligned}$$

- ◆ In normal case, we get $\Delta T < 0$, which means MCR achieves smaller cost (average processing time) than SCR.

Optimization problem of BE class

- ◆ Assumptions:
 - Two classes: EF class and BE class
 - The performance of BE class is mainly concerned because
 - ◆ BE class takes the biggest portion of all network traffic
 - ◆ EF class can be guaranteed due to its high priority.
- ◆ BECOP (BE class optimization problem)
 - Given a DiffServ network $G(V,E)$,
 - ◆ Each pair of nodes is associated with an EF traffic pair with equal rate.
 - ◆ Each pair of nodes is associated with a BE traffic pair with equal rate.
 - Given an EF rate and a routing algorithm r , with the increase of BE rate, the BE rate when some link is saturated is denoted by B_s
 - $B_{\max} = \max\{B_{sr1}, B_{sr2}, \dots\}$ for BE class, when applying different routing algorithms $r1, r2, \dots$
 - BECOP is to find optimal routing algorithm combination of $MCR(r_{ef}, r_{be})$ and B_{\max}

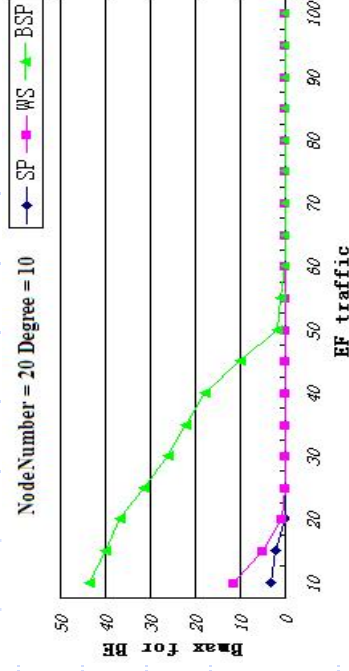
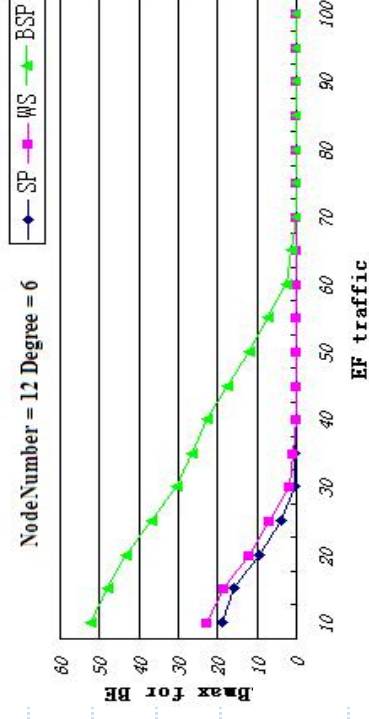
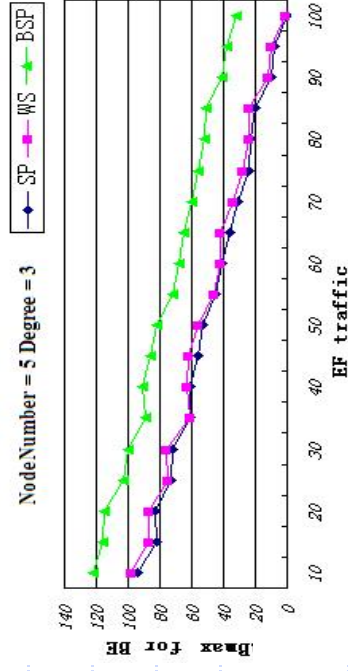
Simulation study

- ◆ Topology generator: thousands of topologies
 - Number of nodes N : 5, 12, 20
 - Variance index for the degree of each node D_{var} : 3, 6, 10
 - Variance index of link capacities C_{var} : 10

Given a specific value of D_{var} we generate a random number within $[1, D_{var}]$ for each node as its degree. And given a specific value of C_{var} we generate a random number within $[100, 100 \times C_{var}]$ for each link as its capacity.

1000 topologies for each configuration
- ◆ Algorithms for EF class and BE class:
 - SP (shortest path)
 - WSP (widest shortest path)
 - BSP (bandwidth-inversion shortest path)

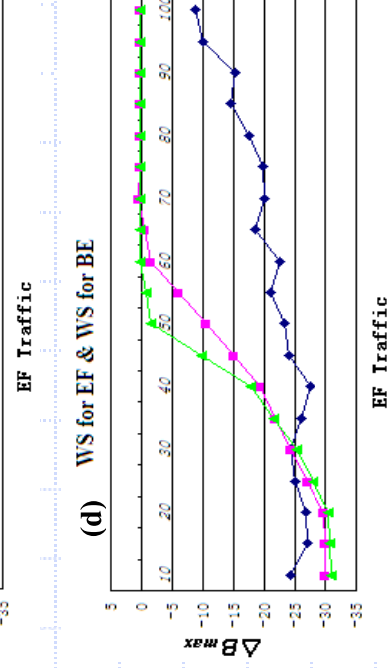
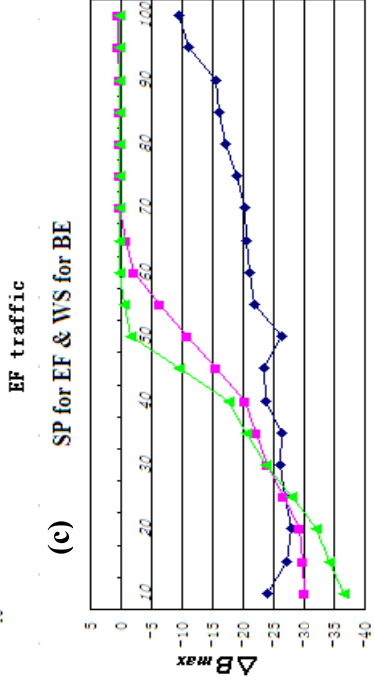
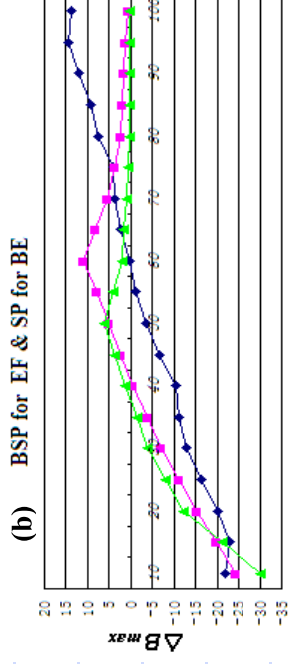
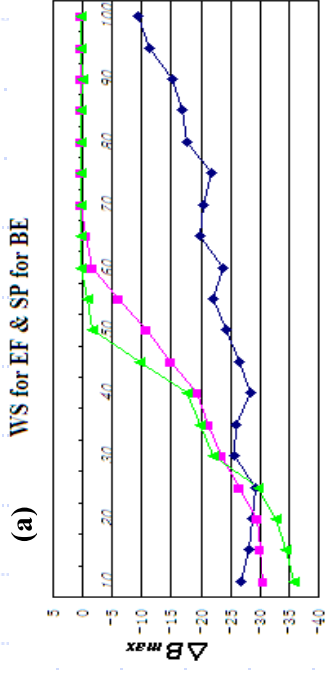
Results – B_{\max} under SCR



- ◆ With the increase of EF traffic, B_{\max} decreases.
- ◆ Given same EF traffic, B_{\max} decreases significantly with the increase of network complexity and size.
- ◆ BSP algorithm achieves the best B_{\max} for BE class under SCR scenario.

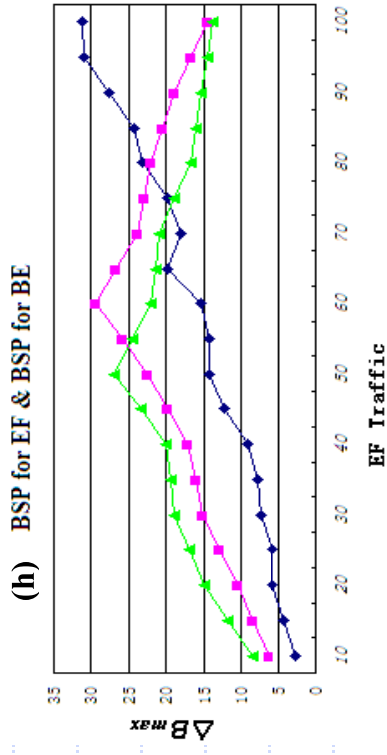
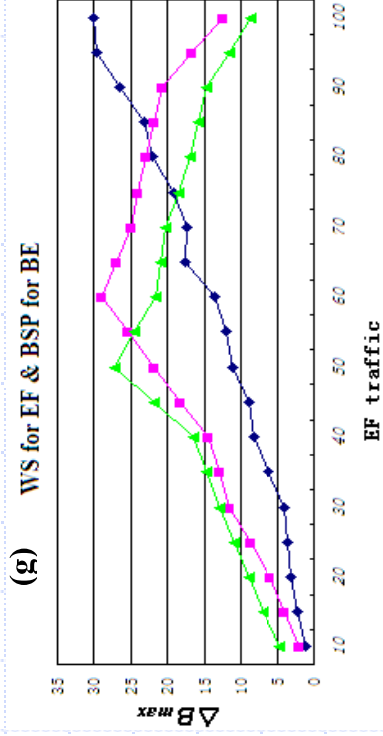
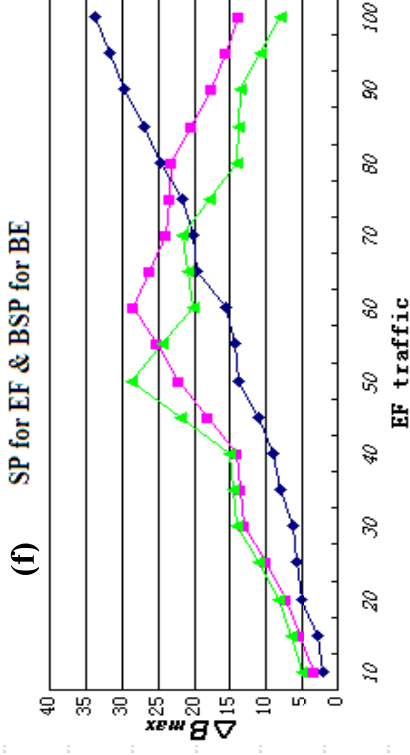
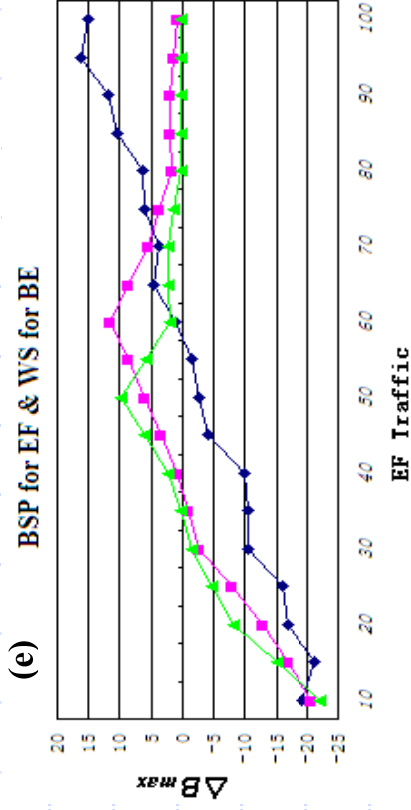
Results – MCR vs. SCR(BSP)

$$\Delta B_{\max} = B_{\max BE}(R_1, R_2) - B_{\max BE}(BSP)$$



◆ These routing combinations of MCR achieve worse performance than SCR(BSP)

Results (cont'd)



- ◆ When BSP is used for BE class, MCR achieves much better performance than SCR(BSP).
- ◆ With the increase of EF traffic, MCR(x, BSP) is becoming better than SCR(BSP)
- ◆ When network becomes more complicated, MCR is much better than SCR.

Prototyping of MCR

- ◆ Based on Zebra open source and major modifications including
 - Single-class routing table → multi-class routing table
 - Normal link state database → TE link state database according to RFC3630
 - Single class route computation → multi-class route computation
- ◆ RFC3630 – Traffic engineering (TE) extension OSPF version 2
 - Became RFC in September 2003
 - Does not specify any QoS routing algorithm, but provides the TE extension to link state database which will become a TE link state database.
 - Has a LSA that can contain different link state information (i.e., available bandwidth) for upto 8 priority classes. This exactly favors the requirement of MCR which needs a link state protocol to exchange link state information for different classes

Future work

- ◆ Experimental study
- ◆ MCR for intra-area study
- ◆ Extend BSP to consider EF traffic ratio on a link

Reference:

Master thesis: Investigation and analysis of optimal multi-class routing algorithms for IP network. June 2003. HUT

Thank you! Any comments?