



Aalto University
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Engineering

Optimal scheduling problem for scalable queues

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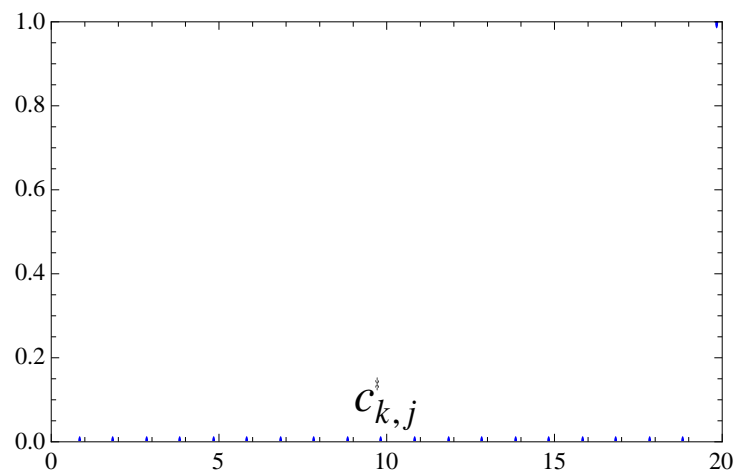
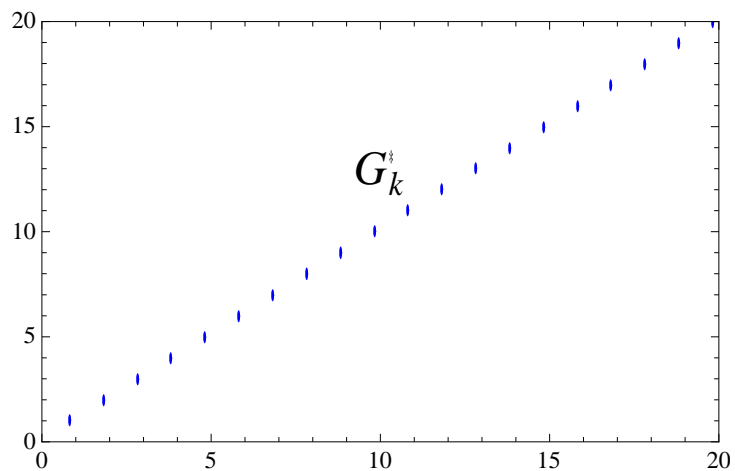
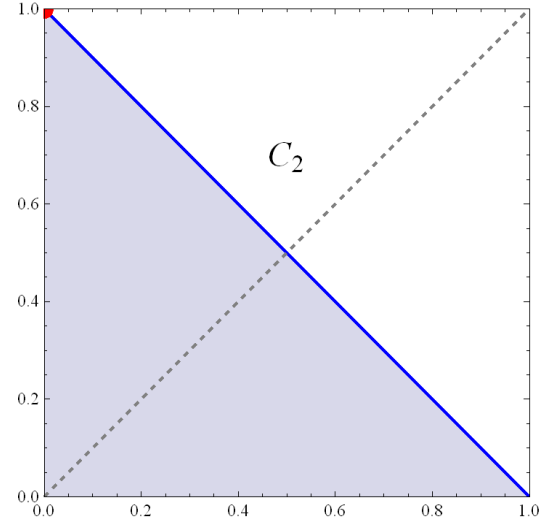
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Dagstuhl Seminar 13111: Scheduling

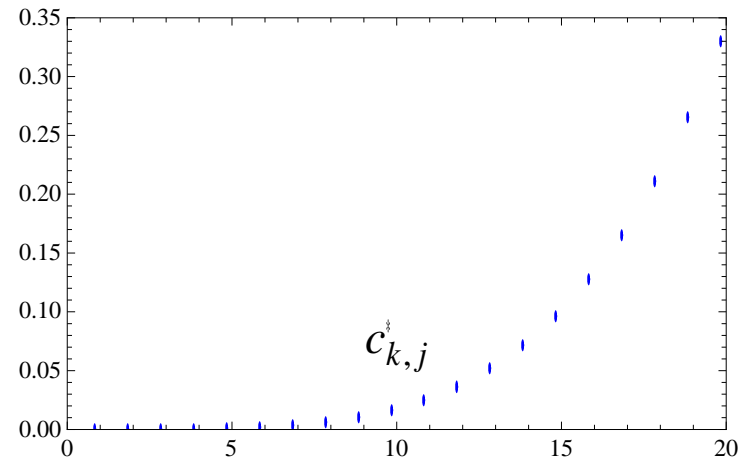
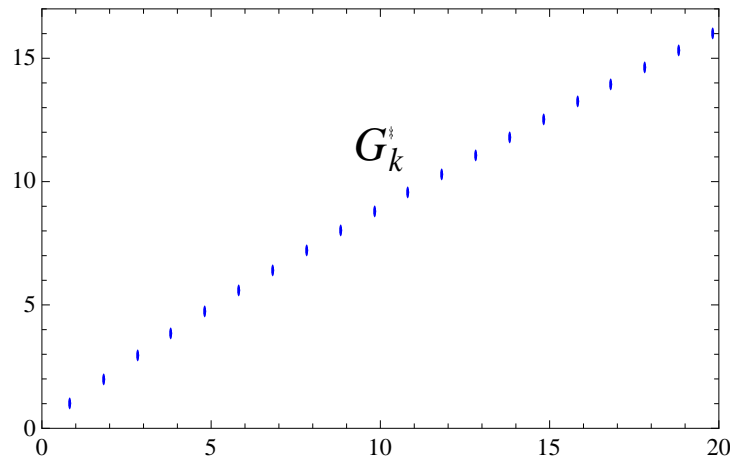
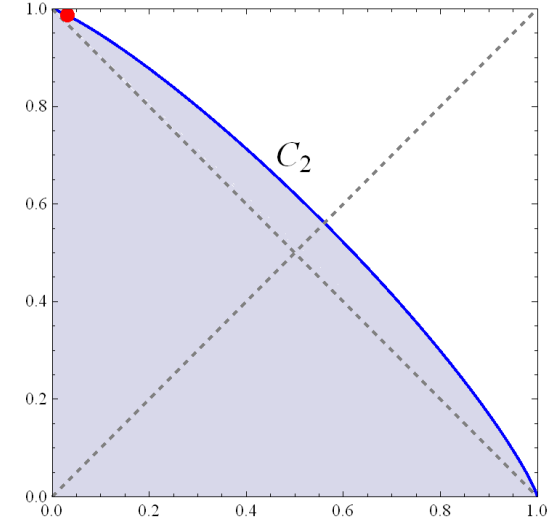
10–15 March 2013

Wadern, Germany

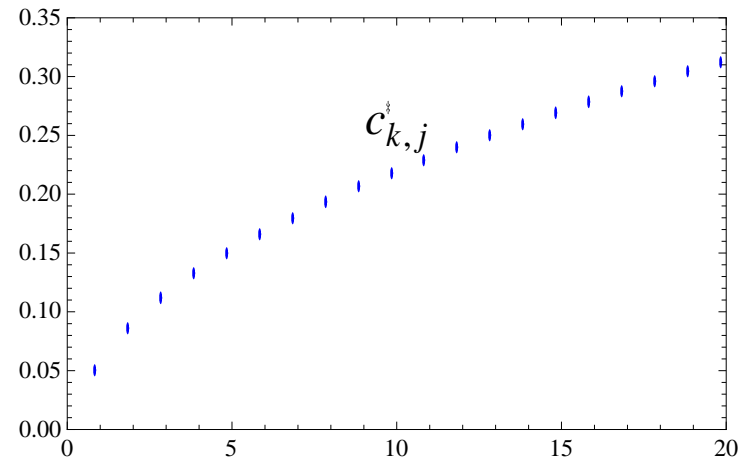
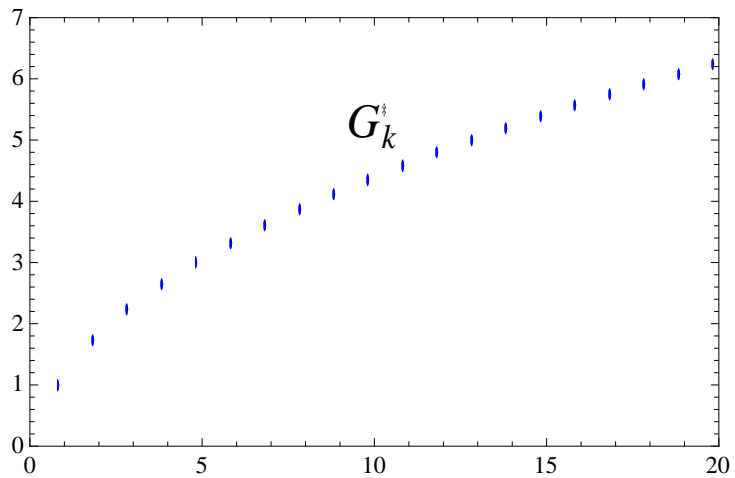
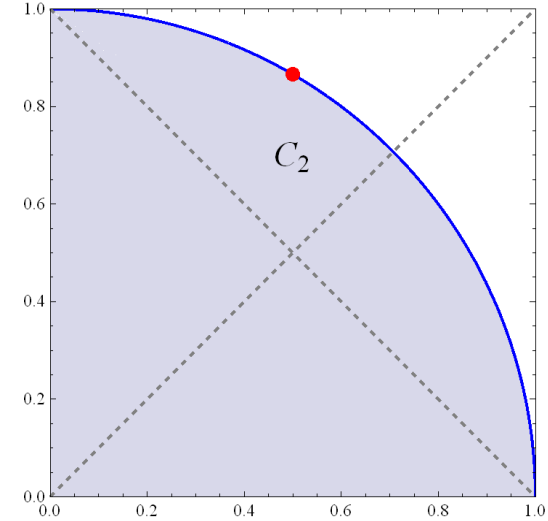
Alpha = 1.0 (single-server queue)



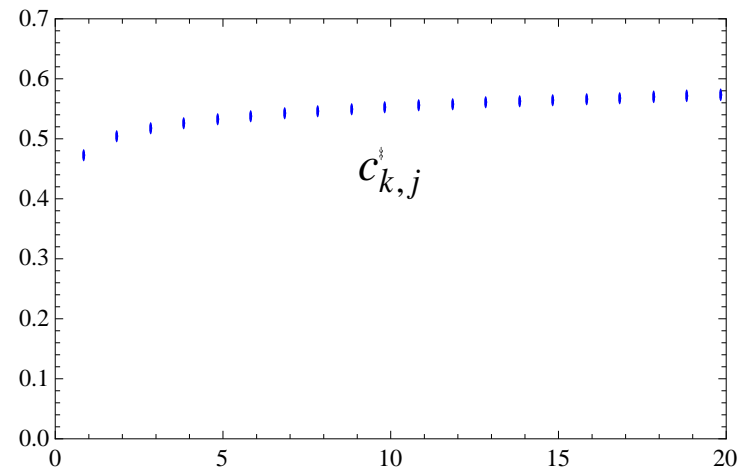
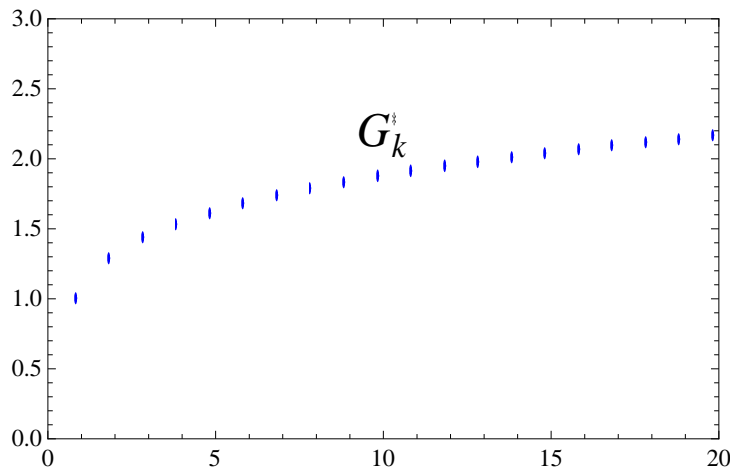
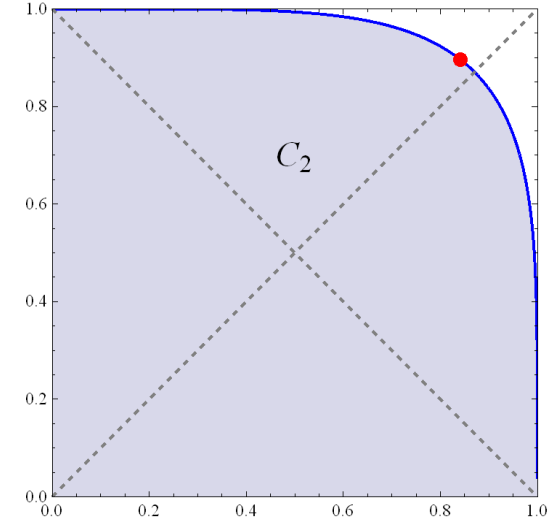
Alpha = 1.2



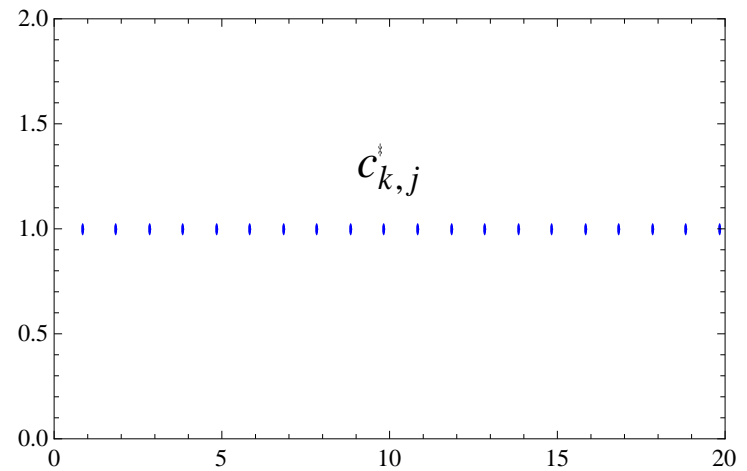
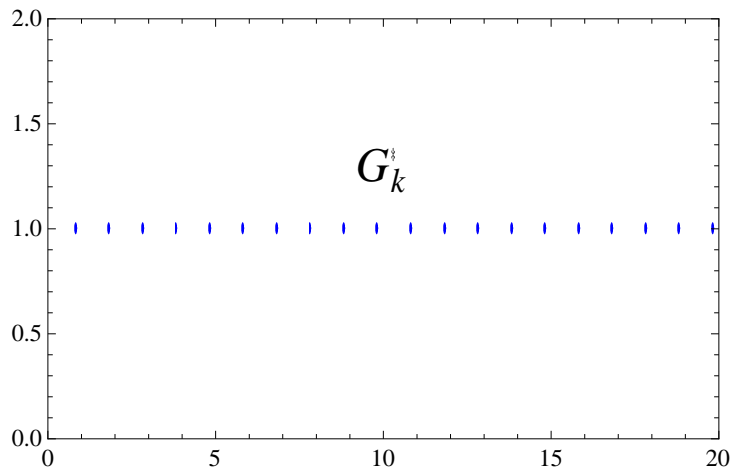
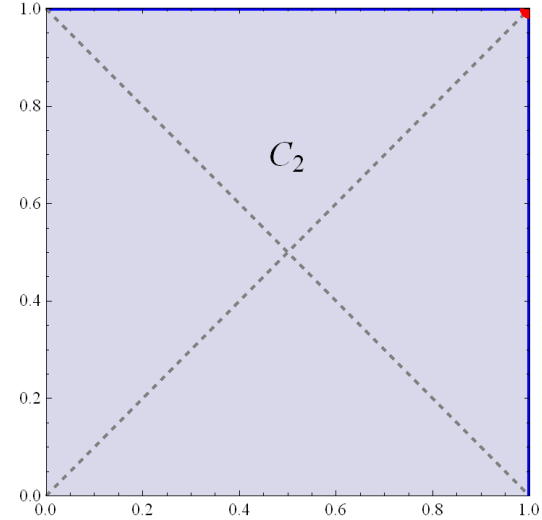
Alpha = 2.0



Alpha = 5.0

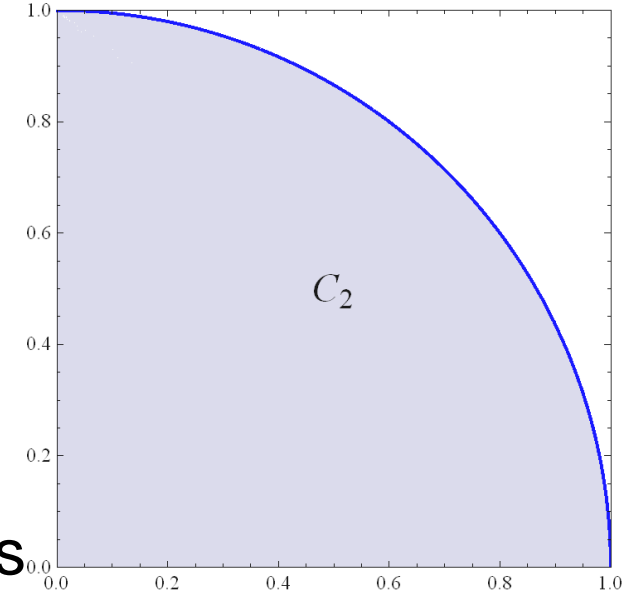


Alpha = infinite (infinite-server queue)



Scalable queue

- Service system where the service capacity scales with number of jobs
- **Policy**: When there are k jobs with sizes



$$s_1 \geq \dots \geq s_k$$

choose a **rate vector**

$$\mathbf{c}_k = (c_{k1}, \dots, c_{kk}) \in C_k$$

and serve job i with rate c_{ki}

- **Assume**: Capacity regions C_k **compact** and **symmetric**

Optimal scheduling problem (transient system without arrivals)

- Assume that there are n jobs in the system at time 0
- What is the optimal way to make the system empty?
- **Objective:** Minimize the mean delay (or flow time)
- **Define:** Flow time for policy ϕ

$$T^\phi = \sum_{i=1}^n t_i^\phi$$

where t_i is the completion time of job i

- **Define:** Operating policies

$$\Phi_n = \{\phi = (\mathbf{c}_1, \dots, \mathbf{c}_n) : \mathbf{c}_k \in C_k \text{ for all } k\}$$

Trivial case: One job

- Define:

$$G_1^* = \frac{1}{c_1^*}, \quad c_1^* = \max_{c_1 \in C_1} c_1$$

- Now

$$T^* = \min_{\phi \in \Phi_1} T^\phi = s_1 G_1^*, \quad \phi^* = (\mathbf{c}_1^*)$$

General case: n jobs

- Define (recursively):

$$G_k^* = \min_{\mathbf{c}_k \in \mathcal{C}_k} g_k(\mathbf{c}_k), \quad g_k(\mathbf{c}_k) = \frac{1}{c_{kk}} \left(k - \sum_{i=1}^{k-1} c_{ki} G_i^* \right)$$

- Theorem [Aalto et al. (2011)]: If

$$G_1^* < \dots < G_n^*$$

then

$$T^* = \min_{\phi \in \Phi_n} T^\phi = \sum_{k=1}^n s_k G_k^*, \quad \phi^* = (\mathbf{c}_1^*, \dots, \mathbf{c}_n^*)$$

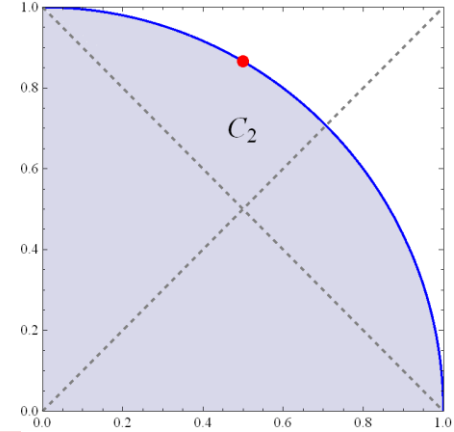
General case: n jobs (cont.)

- In addition,

$$c_{k1}^* \leq \dots \leq c_{kk}^* \text{ for all } k$$

- The optimal policy applies the **SRPT-FM principle**
 - the shortest job is served with the highest rate, etc.
- The optimal rate vector **does not depend on the absolute sizes** of jobs (only on their order)

Alpha-balls



- Let $\alpha \geq 1$ and consider capacity regions

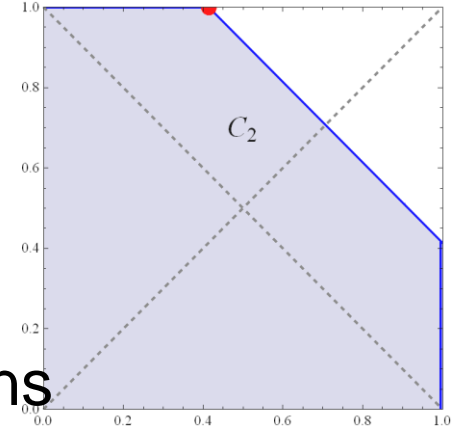
$$C_k = \{\mathbf{c}_k \geq 0 : \sum_{j=1}^k c_{kj}^\alpha \leq 1\}$$

- Now

$$G_k^* = \left(k^{\frac{\alpha}{\alpha-1}} - (k-1)^{\frac{\alpha}{\alpha-1}} \right)^{\frac{\alpha-1}{\alpha}} \quad (\text{increasing in } k)$$

$$c_{kj}^* = \left(\frac{G_j^*}{k} \right)^{\frac{1}{\alpha-1}} \quad (\text{increasing in } j)$$

Symmetric polymatroids



- Let $\gamma_1 < \dots < \gamma_n$ and consider capacity regions

$$C_k = \{\mathbf{c}_k \geq 0 : \sum_{i \in I} c_{ki} \leq \gamma_{|I|}, I \subset \{1, \dots, n\}\}$$

- Theorem:** If $\gamma_1 > \gamma_2 - \gamma_1 > \dots > \gamma_n - \gamma_{n-1}$, then

$$G_1^* < \dots < G_n^* \quad (\text{increasing in } k)$$

$$c_{kj}^* = \gamma_{k-j+1} - \gamma_{k-j} \quad (\text{increasing in } j)$$

- Optimality result of [Sadiq and de Veciana \(2010\)](#)

Open questions

- Is it possible to make the implicit condition explicit?
- Optimal scheduling problem for a dynamic system with random arrivals?
- Other objective functions?

References

- B. Sadiq and G. de Veciana,
Balancing SRPT prioritization vs opportunistic gain in wireless systems with flow dynamics, in *ITC-22, 2010*
- S. Aalto, A. Penttinen, P. Lassila and P. Osti,
On the optimal trade-off between SRPT and opportunistic scheduling, in *ACM SIGMETRICS, 2011*
- S. Aalto, A. Penttinen, P. Lassila and P. Osti,
Optimal size-based opportunistic scheduler for wireless systems, *Queueing Systems 72, 2012*

The End